

**Pre A\*- Clones, “If-Then-Else” Structures over Pre A\*- Algebra**<sup>1</sup>K. Suguna Rao and <sup>2</sup>P.Koteswara Rao<sup>1</sup>Dept. of Mathematics, Acharya Nagarjuna University, Nagarjuna Nagar-522 510,  
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**Abstract:** In this paper we divide two sections. In the first section we introduce the Definition of Pre A\*-Clone, Every Pre A\*-Clone generates Pre A\*-algebra, Pre A\*-algebra forms Pre A\*-Clone and some other theorems. Definition of Boolean Clone ,Every Boolean clone generates Boolean Algebra. In the Second Section ,we introduce the definition of if-then –else operation on pre A\*-algebra and pointed set and some other theorems.

**Key Words:** Clone, A\*-Clone, Boolean Algebra, Huntington’s theorem.

**1. Pre A\*- Clone:**

**Definition:** A Pre A\*- Clone over a set ‘S’ is a set A of ternary operations on S satisfying the following conditions.

$$\forall f, g \in A, x, y, z, x_1, y_1, z_1 \in S$$

$$(i) f(x, x, x) = x$$

$$(ii) f(f(x_1, y_1, z_1), f(x_2, y_2, z_2), f(x_3, y_3, z_3)) = f(x_1, y_2, z_3)$$

$$(iii) f(g(x_1, y_1, z_1), g(x_2, y_2, z_2), g(x_3, y_3, z_3)) = g(f(x_1, x_2, x_3), f(y_1, y_2, y_3), f(z_1, z_2, z_3))$$

$$(iv) A \text{ contains the projections } \Pi_1, \Pi_2 \text{ where } \Pi_1(x, y, z) = x, \Pi_2(x, y, z) = y$$

Further ‘A’ is closed under the operations  $f \uparrow, f \wedge g, f \vee g$  defined by

$$(v) f \uparrow(x, y, z) = f(y, x, z)$$

$$(vi) (f \wedge g)(x, y, z) = f(g(x, y, z), g(y, y, z), z)$$

$$(vii) (f \vee g)(x, y, z) = f(g(x, x, z), g(x, y, z), z)$$

**1.1 Example:** Let  $X, Y$  be sets and 'S' be the set of all partial functions from  $X$  into  $Y$ . Then  $3^X$  is a pre  $A^*$ -Clone over  $S$  w. r. t to the operations defined by

$$a(x, y, z)(P) = \begin{cases} x(P) & \text{if } a(P) = 1 \\ y(P) & \text{if } a(P) = 0 \\ z(P) & \text{if } a(P) = 2 \end{cases}$$

For  $a \in 3^X, x, y, z, \in S$  and  $P$  lies in the domains of  $x, y, z$ .

**1.2 Theorem:** Every a pre  $A^*$ -Clone on a set 'S' is a pre  $A^*$ -Algebra.

**Proof:** Suppose  $A$  is a pre  $A^*$ -Clone over a set 'S'. Clearly  $f \wedge g = g \wedge f$

$$(f \vee g)(x, y, z) = f(g(x, x, z), g(x, y, z), z)$$

(i)  $\underline{f \sim = f, \forall f \in A}$

$$\begin{aligned} f \sim(x, y, z) &= (f \sim)(x, y, z) \\ &= f \sim(y, x, z) \\ &= f(x, y, z) \\ &= f \end{aligned}$$

$$\therefore f \sim = f$$

(ii)  $\underline{f \wedge f = f}$

$$\begin{aligned} (f \wedge f)(x, y, z) &= f(f(x, y, z), f(y, y, z), z) \\ &= f(x, y, z), f(y, y, z), z \\ &= f(x, y, z). \end{aligned}$$

$$\therefore f \wedge f = f$$

(iii)  $\underline{f \wedge g = g \wedge f}$

$$\begin{aligned} (f \wedge g)(x, y, z) &= f(g(x, y, z), g(y, y, z), z) \\ &= g(f(x, y, z), f(y, y, z), f(z, z, z)) \\ (g \wedge f)(x, y, z) &= g(f(x, y, z), f(y, y, z), z) \\ &= g(f(x, y, z), f(y, y, z), f(z, z, z)) \end{aligned}$$

$$\therefore f \wedge g = g \wedge f$$

(iv)  $\underline{(f \wedge g) \sim = f \sim \vee g \sim}$

$$\begin{aligned} (f \wedge g) \sim(x, y, z) &= (f \wedge g)(y, x, z) \\ &= f(g(y, x, z), g(x, x, z), z) \\ (f \sim \vee g \sim)(x, y, z) &= f \sim(g(x, y, z), g(x, y, z), z) \end{aligned}$$

$$= \tilde{f}(g(x, y, z), g(x, y, z), z)$$

$$= f(g(y, x, z), g(x, x, z), z)$$

$$\therefore (f \wedge g)^\sim = \tilde{f} \vee g^\sim$$

$$(v) \underline{f \wedge (g \wedge h) = (f \wedge g) \wedge h}$$

$$(f \wedge (g \wedge h))(x, y, z) = f(g \wedge h(x, y, z))$$

$$= f(g(h(x, y, z), h(y, y, z), z))$$

$$= (f \wedge g)(h(x, y, z), h(y, y, z), z)$$

$$= (f \wedge g)(h(x, y, z))$$

$$= ((f \wedge g) \wedge h)(x, y, z)$$

$$\therefore f \wedge (g \wedge h) = (f \wedge g) \wedge h$$

$$(vi) \underline{f \wedge (g \vee h) = (f \wedge g) \vee (f \wedge h)}$$

$$\text{Now } (f \wedge (g \vee h))(x, y, z) = f((g \vee h)(x, y, z))$$

$$= f(g(h(x, x, z), h(x, y, z), z))$$

$$= (f \wedge g)(f(h(x, x, z), h(x, y, z), z))$$

$$= (f \wedge g) \wedge f(h(x, y, z))$$

$$= (f \wedge g) \vee (f \wedge h)$$

$$\text{Now } (f \wedge g)(x, y, z) = f(g(x, y, z), g(y, y, z), z)$$

$$(f \wedge h)(x, y, z) = f(h(x, y, z), h(y, y, z), z)$$

$$\therefore ((f \wedge g) \vee (f \wedge h))(x, y, z) = f(g(x, y, z), g(y, y, z), z) \vee f(h(x, y, z), h(y, y, z), z)$$

$$= f((g(x, y, z), g(y, y, z), z) \vee (h(x, y, z), h(y, y, z), z))$$

$$(vii) \quad x \wedge y = x \wedge (x^\sim \wedge y) \text{ for all } x$$

$$f \wedge g = f \wedge (f^\sim \wedge g) \text{ for all } f, g \in A$$

$$f \wedge g(x, y, z) = f(g(x, y, z), g(y, y, z), z)$$

$$f^\sim \wedge g(x, y, z) = \tilde{f}(g(x, y, z), g(y, y, z), z)$$

$$(f \wedge (f^\sim \wedge g))(x, y, z) = f(f(g(y, y, z), g(x, y, z), z))$$

$$\therefore A \text{ is a pre } A^* \text{-Algebra.}$$

**1.3Theorem:** Every pre A\*-Algebra is a pre A\*-Clone over a set 'S'.

**Proof:** Suppose  $(A, \wedge, \vee, (-)^\sim)$  A is a pre A\*-Algebra

Let  $a \in A, x, y, z \in B(A)$  define  $a(x, y, z) = (a_x) \vee (a_y) \vee (a_z)$  for all  $x, y, z \in B(A)$

Then A is a set of ternary operations on B(A).

$$(i) a(x, x, x) = (ax) \vee (ax) \vee (ax)$$

$$= (a \vee a \vee a)(x)$$

$$= 1. x = x$$

$$\therefore a(x, x, x) = x$$

$$(ii) a(a(x_1, y_1, z_1), a(x_2, y_2, z_2), a(x_3, y_3, z_3))$$

$$= a((ax_1) \vee (ay_1) \vee (az_1)) \vee a((ax_2) \vee (ay_2) \vee (az_2)) \vee a((ax_3) \vee (ay_3) \vee (az_3))$$

$$= ax_1 \vee ay_2 \vee az_3$$

$$= a(x_1, y_2, z_3)$$

$$\therefore a(a(x_1, y_1, z_1), a(x_2, y_2, z_2), a(x_3, y_3, z_3)) = a(x_1, y_2, z_3)$$

$$(iii) a(b(x_1, y_1, z_1), b(x_2, y_2, z_2), b(x_3, y_3, z_3))$$

$$= a(bx_1 \vee by_1 \vee bz_1) \vee a(bx_2 \vee by_2 \vee bz_2) \vee a(bx_3 \vee by_3 \vee bz_3)$$

$$= b(ax_1 \vee ax_2 \vee ax_3) \vee b(ay_1 \vee ay_2 \vee ay_3) \vee b(az_1 \vee az_2 \vee az_3)$$

$$= b(a(x_1 \vee x_2 \vee x_3), a(y_1 \vee y_2 \vee y_3), a(z_1, z_2, z_3))$$

$$\therefore a(b(x_1, y_1, z_1), b(x_2, y_2, z_2), b(x_3, y_3, z_3)) = b(a(x_1 \vee x_2 \vee x_3), a(y_1 \vee y_2 \vee y_3), a(z_1, z_2, z_3))$$

$$(iv) \quad \Pi_1 = 1, \quad \Pi_2 = 0, \text{ since}$$

$$\Pi_1(x, y, z) = 1(x, y, z)$$

$$= 1x \vee 1y \vee 1z$$

$$= x$$

$$\parallel y, \Pi_2(x, y, z) = y$$

$$(v) a(b(x, y, z), b(y, y, z), z)$$

$$= (abx) \vee (aby) \vee (abz) \vee (aby) \vee (aby) \vee (abz) \vee (az)$$

$$\begin{aligned}
&= (abx) \vee (aby) \vee (aby) \vee (aby) \vee (a \vee a)bz \vee az \\
&= (abx) \vee (aby) \vee (aby) \vee (aby) \vee (ab \vee a)z \\
&= (abx) \vee (aby) \vee (aby) \vee (aby) \vee (aby) \vee (a \wedge b)z \\
&= (a \wedge b)x \vee (a \wedge b)y \vee (a \wedge b)z \\
&= (a \wedge b)(x, y, z)
\end{aligned}$$

$$\therefore (a \wedge b)(x, y, z) = a(b(x, y, z), b(y, y, z), z)$$

$$(vi) a(b(x, x, z), b(x, y, z), z)$$

$$\begin{aligned}
&= (abx) \vee (abx) \vee (abz) \vee (abx) \vee (aby) \vee (abz) \vee (az) \\
&= (abx) \vee (abx) \vee (abx) \vee (aby) \vee (a \vee a)bz \vee az \\
&= (abx) \vee (abx) \vee (abx) \vee (aby) \vee (ab \vee a)z \\
&= (ab \vee ab \vee ab \vee a)z \\
&= (a \wedge b)(x, y, z)
\end{aligned}$$

$$(vii) a^{\sim}(x, y, z) = (a^{\sim}x) \vee (a^{\sim}y) \vee (a^{\sim}z)$$

$$= a^{\sim}y \vee a^{\sim}x \vee a^{\sim}z$$

$$= ay \vee ax \vee az$$

$$= a(y, x, z)$$

$$\therefore a^{\sim}(x, y, z) = a(y, x, z)$$

$\therefore A$  is a Pre  $A^*$ -clone over  $B(A)$

**1.4 Definition:** A set  $B$  of binary operations on a set  $S$  satisfying the following conditions is called a Boolean Clone over 'S'.

$$(i) f(x, x) = x, \forall x \in S, \forall f \in B$$

$$(ii) f(f(x, y), f(u, v)) = f(x, v)$$

$$(iii) B \text{ contains projections } \pi_1, \pi_2 \text{ where } \pi_1(x, y) = x, \pi_2(x, y) = y$$

$$(iv) f(g(x, y), g(u, v)) = g(f(x, u), f(y, v)) \text{ further } B \text{ is closed under the operations } f^l, f \wedge g \text{ defined by}$$

$$(v) f^l(x, y) = f(y, x) \quad (vi) (f \wedge g)(x, y) = f(g(x, y), y)$$

**1.5 Example:** For any let  $X$ ,  $2^X$  is a Boolean clone over the Set 'S' of all functions from  $X$  into a given set  $Y$  where we define for  $f \in 2^X$ ,  $x, y \in S$ ,  $P \in X$ .

$$\begin{aligned} f(x, y)(P) &= x(P) \text{ if } f(P) = 1 \\ &= y(P) \text{ if } f(P) = 0 \end{aligned}$$

**1.6 Theorem:** Every Boolean Clone  $B$  over a Set 'S' becomes a Boolean algebra.

$$\begin{aligned} (f \wedge f^l)(x, y) &= f(f^l(x, y), y) \\ &= f(f(y, x), y) \\ &= f(f(y, x), f(y, y)) \\ &= f(y, y) = y = \Pi_2(x, y) \end{aligned}$$

$$\therefore (f \wedge f^l) = \Pi_2$$

$$\begin{aligned} [(f \wedge g) \vee (f \wedge g^l)](x, y) &= (f \wedge g)[x, f \wedge g^l(x, y)] \\ &= (f \wedge g)[x, f(g(y, x), y)] \\ &= (f \wedge g)[x, g(y, f(x, f))] \\ &= f[g(x, (g(y, f(x, y)))), g(y, f(x, y))] \\ &= f(g(x, f(x, y)), g(y, f(x, y))) \\ &= g(f(x, y), f(f(x, y), f(x, y))) \\ &= g(f(x, y), f(x, y)) \\ &= f(x, y) \end{aligned}$$

$$\therefore [(f \wedge g) \vee (f \wedge g^l)] = f$$

By Huntington's theorem,  $(B, \wedge, \vee, (-)^l, \Pi_2)$  is a Boolean Algebra.

**1.7 Corollary:** If  $A$  is Pre  $A^*$ -clone over 'S', then every Sun pre  $A^*$ -Algebra of  $A$  is a Pre  $A^*$ -clone over  $S$ .

**1.8 Theorem:** If  $A$  is Pre  $A^*$ -clone over  $S$  then there exists a set  $X$  such that  $A$  is isomorphic to a pre sub clone of  $3^X$ .

**Proof:** Since A is a Pre A\*-Algebra

⇒ A is isomorphic to a sub Pre A\*-Algebra  $P F_n(x, 2)$  for some X and since  $P F_n(x, 2)$  can be identified with  $3^x$ , A is isomorphic to sub Pre A\*-Algebra of  $3^x$ . Let Y be any set and S be the set of all partial functions from X to Y. Then  $3^x$  is a Pre A\*-clone over S. Since every sub Pre A\*-Algebra of  $3^x$  is an A\*-clone over S, A is isomorphic to a sub Pre A\*-clone of  $3^x$ .

## 2 If-Then-Else Operation:

**2.1 Definition:** Let 'A' be a pre A\* -algebra with 1. If  $x a_1 b \in A$ ; define the If - then - else operation on A as  $if_x(a, b) = (x \wedge a) \vee (x \sim \wedge b)$ .

( $if_x(a, b) =$  should be viewed as conditional "if  $x_1$  then  $a_1$  else b")

**2.2 Lemma:** Every pre A\*-algebra which indicates Constants satisfies the following laws.

- (i)  $if_2(a_1 b) = 2$
- (ii)  $if_x(2, 2) = 2$
- (iii)  $if_1(a_1 b) = a(b \neq 2)$
- (iv)  $if_0(a_1 b) = b(p \neq 2)$
- (v)  $if_x(1, 0) = x$

**Proof:**

- (i)  $if_2(a, b) = (2 \wedge a) \vee (2 \sim \wedge b) = 2 \vee (2 \wedge b) = 2 \vee 2 = 2$
- (ii)  $if_x(2, 2) = (x \wedge 2) \vee (x \sim \wedge 2) = 2 \vee 2 = 2$
- (iii)  $if_1(a, b) = (1 \wedge a) \vee (1 \sim \wedge b) = a \vee (0 \wedge b) = a \vee 0 (b \neq 2) = a$
- (iv)  $if_0(a, b) = (0 \wedge a) \vee (0 \sim \wedge b) = 0 \vee (1 \wedge b)(p \neq 2) = 0 \vee b = b$
- (v) If  $x = 2$ , then  $if_x(1, 0) = (2 \wedge 1) \vee (2 \sim \wedge 0) = 2 = x$   
 If  $x \neq 2$ , then  $if_x(1, 0) = (x \wedge 1) \vee (x \sim \wedge 0) = x \vee 0 = x$

One can observe that either x or a or b is 2 then  $if_x(a, b) = 2$

**2.3 Lemma:** Every pre A\*-algebra satisfies

- (i)  $if_{x \sim}(a, b) = if_x(b, a)$
- (ii)  $if_x(a, b) \wedge c = if_x(a \wedge c), (b \wedge c)$
- (iii)  $if_x(a, b) \vee c = if_x(a \vee c), (b \vee c)$
- (iv)  $if_x(if_y(a, b), if_y(c, d)) = if_y(if_x(a, c), if_x(b, d))$
- (v)  $if_{x \vee y}(a, b) = if_x(a, if_y(a, b))$

- (vi)  $if_{x \wedge y}(a, b) = if_x(if_y(a, b), b)$   
 (vii)  $if_p(a, a) = a$

**Proof:**

(i)  $if_{x \sim}(a, b) = (x \sim \wedge a) \vee (x \sim \sim \wedge b)$   
 $= (x \sim \wedge a) \vee (x \wedge b)$   
 $= (x \wedge b) \vee (x \sim \wedge a)$   
 $= if_x(b, a)$

Therefore  $if_{x \sim}(a, b) = if_x(b, a)$

(ii)  $if_x(a, b) \wedge c = \{(x \wedge a) \vee (x \sim \wedge b)\} \wedge c$   
 $= \{(x \wedge a) \wedge c\} \vee \{(x \sim \wedge b) \wedge c\}$   
 $= \{x(a \wedge c)\} \vee \{x \sim (b \wedge c)\}$   
 $= if_x(a \wedge c, b \wedge c)$

(iii) If  $x \neq 2$ , then

$$if_x(a \vee c, b \vee c) = \{x \wedge (a \vee c)\} \vee \{x \sim \wedge (b \vee c)\}$$

$$= \{(x \wedge a) \vee (x \wedge c)\} \vee \{(x \sim \wedge b) \vee (x \sim \wedge c)\}$$

$$= \{(x \wedge a) \vee (x \sim \wedge b)\} \vee \{(x \wedge c) \vee (x \sim \wedge c)\}$$

$$= \{(x \wedge a) \vee (x \sim \wedge b)\} \vee \{(x \vee x \sim) \wedge c\}$$

$$= if_x(a, b) \vee (1 \wedge c) \quad (x \neq 2 \Rightarrow x \vee x \sim = 1)$$

$$= if_x(a, b) \vee c$$

If  $x = 2$ , the proof is clear.

(iv)  $if_x(if_y(a, b), if_y(c, d)) = \{x \wedge if_y(a, b)\} \vee \{x \sim \wedge if_y(c, d)\}$   
 $= \{x \wedge [(y \wedge a) \vee (y \sim \wedge b)]\} \vee \{x \sim \wedge [(y \wedge c) \vee (y \sim \wedge d)]\}$   
 $= \{x \wedge (y \wedge a)\} \vee \{x \wedge (y \sim \wedge b)\} \vee \{x \sim \wedge (y \wedge c)\} \vee \{x \sim \wedge (y \sim \wedge d)\}$   
 $= \{y \wedge (x \wedge a)\} \vee \{y \wedge (x \sim \wedge c)\} \vee \{y \sim \wedge (x \wedge b)\} \vee \{y \sim \wedge (x \sim \wedge d)\}$   
 $= [y \wedge \{(x \wedge a) \wedge (x \sim \wedge c)\}] \vee [y \sim \wedge \{(x \wedge b) \wedge (x \sim \wedge d)\}]$   
 $= [y \wedge if_x(a, c)] \vee [y \sim \wedge if_x(b, d)]$   
 $= if_y(if_x(a, c), if_x(b, d))$

(v)  $if_x(a, if_y(a, b)) = (x \wedge a) \vee (x \sim \wedge if_y(a, b))$   
 $= (x \wedge a) \vee [x \sim \wedge \{(y \wedge a) \vee (y \sim \wedge b)\}]$   
 $= (x \wedge a) \vee [\{x \sim \wedge (y \wedge a)\} \vee \{x \sim \wedge (y \sim \wedge b)\}]$   
 $= [(x \wedge a) \vee \{x \sim \wedge (y \wedge a)\}] \vee \{x \sim \wedge (y \sim \wedge b)\}$



$$\begin{aligned}
 &= [(x \wedge a) \vee \{(x \sim \wedge y) \wedge a\}] \vee \{x \sim \wedge (y \sim \wedge b)\} \\
 &= [\{x \vee \{(x \sim \wedge y)\} \wedge a\} \vee \{x \sim \wedge (y \sim \wedge b)\}] \\
 &= \{(x \vee y) \wedge a\} \vee \{(x \vee y) \sim \wedge b\} \\
 &= \text{if}_{x \vee y}(a, b)
 \end{aligned}$$

$$\begin{aligned}
 \text{(vi)} \quad \text{if}_x(\text{if}_y(a, b), b) &= (x \wedge \text{if}_y(a, b)) \vee (x \sim \wedge b) \\
 &= [x \wedge \{(y \wedge a) \vee (y \sim \wedge b)\}] \vee (x \sim \wedge b) \\
 &= [(x \wedge (y \wedge a)) \vee (x \wedge (y \sim \wedge b))] \vee (x \sim \wedge b) \\
 &= (x \wedge (y \wedge a)) \vee [(x \wedge (y \sim \wedge b)) \vee (x \sim \wedge b)] \\
 &= (x \wedge (y \wedge a)) \vee [(x \wedge y \sim) \wedge b] \vee (x \sim \wedge b) \\
 &= (x \wedge (y \wedge a)) \vee [\{(x \wedge y \sim) \vee x \sim\} \wedge b] \\
 &= \{(x \wedge y) \wedge a\} \vee \{(x \sim \vee y \sim) \wedge b\}. \\
 &= \{(x \wedge y) \wedge a\} \vee \{(x \wedge y) \sim \wedge b\}. \\
 &= \text{if}_{x \wedge y}(a, b)
 \end{aligned}$$

$$\begin{aligned}
 \text{(vii)} \quad \text{if}_p(a, a) &= (a \wedge a) \vee (a \sim \wedge a) \\
 &= a \vee (a \sim \wedge a) \\
 &= a
 \end{aligned}$$

$$\begin{aligned}
 \text{Observe that } \text{if}_x(a, a) &= a(x \wedge a) \vee (x \sim \wedge a) \\
 &= (x \vee x \sim) \wedge a, \text{ if } x \neq 2,
 \end{aligned}$$

then  $x \vee x \sim = 1$  and  $\text{if}_x(a, a) = a$  Otherwise the value is 2.

**2.4 Definition:** A pointed set  $(S, \perp)$  it is said to be an “if-then-else” structure over pre  $A^*$ -Algebra  $A$ ; if the ternary operation

$$\begin{aligned}
 (x, a, y) &\rightarrow xay, \forall s, y \in S, \forall a \in A \\
 SXAXS &\rightarrow S \text{ satisfied}
 \end{aligned}$$

- (i)  $xax = xa \sim \perp$
- (ii)  $x \perp y = x$
- (iii)  $xa \sim y = yax$
- (iv)  $(xay)bz = xab[(\perp bz)a(ybz)]$  where  $ab = a \wedge b$

**2.5 Note:**  $a \sim = (a \vee 1)$  is called a hault.

**2.6 Lemma:** If  $(S, \perp)$  is an “if-then-else” structure over pre  $A^*$ -Algebra  $A$ , then

- (i)  $xay = xa(\perp ay)$
- (ii)  $xab[(\perp bz)a(ybz)] = (\perp bz)a(ybz)$

$$(iii) (xby)a(\mu b\theta) = (xa\mu)b(ya\theta)$$

$$(iv) x(a \wedge b)y = (xay)b(yay)$$

**Proof:**

$$(i) \quad xay = x(a \vee a)y \\ = xa(\perp ay)$$

$$\therefore xay = xa(\perp ay)$$

$$(ii) \quad xab[(\perp bz)a(ybz)]$$

$$= (\perp bz)a[xab(yba(\perp bz))]$$

$$[\because xa(ybz) = yb(xaz)] \quad \text{where } ab = a \wedge b$$

$$= (\perp bz)a[\perp b(xab(ybz))]$$

$$= (\perp bz)a[\perp b(xay)b(zaz)]$$

$$= (\perp bz)a[\perp b\{(xa(yay))b(za(zaz))\}]$$

$$= (\perp bz)a[\perp b\{(xbz)a(ybz)a(ybz)\}]$$

$$= (\perp bz)a[\perp b\{(xbz)a(yay)b(zaz)\}]$$

$$= (\perp bz)a[(\perp a \perp)b\{(xbz)a[(ybz)a(ybz)]\}]$$

$$= (\perp bz)a[(xbz)a(ybz)]$$

$$\therefore xab[(\perp bz)a(ybz)] = (\perp bz)a[(xbz)a(ybz)]$$

$$(iii) \quad (xby)a(\mu b\theta)$$

$$= [xb(\perp by)]a[\mu b(\perp b\theta)]$$

$$= [xb(\perp by)]a[\perp a[\mu b(\perp b\theta)]]$$

$$= [xb(\perp by)]a\{[\perp b(\perp b \perp)]a[\mu b(\perp b\theta)]\}$$

$$= [xb(\perp by)]a\{(\perp a\mu)b[(\perp b \perp)a(\perp b\theta)]\}$$

$$= [xb(\perp by)]a\{(\perp a\mu)b[(\perp a \perp)b(\perp a\theta)]\}$$

$$= [xa(\perp a\mu)]b\{(\perp by)a[(\perp a \perp)b(\perp a\theta)]\}$$

$$= [xa(\perp a\mu)]b\{[\perp a[(\perp a \perp)]b[ya(\perp a\theta)]\}$$

$$= (xa\mu)b[\perp b[ya(\perp a\theta)]]$$

$$= (xa\mu)b[\perp b(ya\theta)]$$

$$= (xa\mu)b(ya\theta)$$

$$\therefore (xby)a(\mu b\theta) = (xa\mu)b(ya\theta)$$

$$(iv) \quad (xay)b(yay)$$

$$= [\perp a(xay)]b[\perp a(yay)]$$

$$= [\perp a(xay)]b[\perp b(\perp ay)]$$

$$\begin{aligned}
 &= \perp b\{\perp a(xay)\}b(\perp ay) \\
 &= \perp b\{xa(\perp ay)\}b(\perp ay) \\
 &= \perp b\{xab(\perp ay)\} \\
 &= xab(\perp b(\perp ay)) \\
 &= xab(\perp a \vee by) \\
 &= x(a \wedge b)y \\
 &\therefore x(a \wedge b)y = (xay)b(yay)
 \end{aligned}$$

**2.7 Theorem:** pre A\*-Algebra A is a pre A\*-clone iff  $(S, \perp)$  is an “if-then-else” structure over A for some distinguished element  $\perp$ .

**Proof:** Case-I: Suppose  $(S, \perp)$  is an “if-then-else” structure over A.

For  $x, y, z \in S, a \in A$ , define

$$a(x, y, z) = za(xay)$$

For  $a(x, x, x) = x$  is clear

$$\begin{aligned}
 &a(a(x_1, y_1, z_1), a(x_2, y_2, z_2), a(x_3, y_3, z_3)) \\
 &= [z_3 a(x_3 a y_3)] a\{[z_1 a(x_1 a y_1)]a[z_2 a(x_2 a y_2)]\} \\
 &= [z_3 a\{[z_1 a(x_1 a y_1)]a[z_2 a(x_2 a y_2)]\}] \\
 &= [z_3 a\{(x_1 a y_1) a[z_2 a(x_2 a y_2)]\}] \\
 &= z_3 a\{x_1 a[z_2 a(x_2 a y_2)]\} \\
 &= z_3 a\{z_2 a[x_1 a(x_2 a y_2)]\} \\
 &= z_3 a\{z_2 a[x_1 a y_2]\} \\
 &= z_3 a(x_1 a y_2) = a(x_1, y_2, z_3)
 \end{aligned}$$

$$\therefore a(a(x_1, y_1, z_1), a(x_2, y_2, z_2), a(x_3, y_3, z_3)) = a(x_1, y_2, z_3)$$

$$\begin{aligned}
 &a(b(x_1, y_1, z_1), b(x_2, y_2, z_2), b(x_3, y_3, z_3)) \\
 &= a[z_1 b(x_1 b y_1), z_2 b(x_2 b y_2), z_3 b(x_3 b y_3)] \\
 &= [z_3 b(x_3 b y_3)]a\{[z_1 b(x_1 b y_1)] a [z_2 b(x_2 b y_2)]\} \\
 &= [z_3 b(x_3 b y_3)]a\{(z_1 a z_2) b[(x_1 b y_1) a (x_2 b y_2)]\} \\
 &= [z_3 b(x_3 b y_3)]a\{(z_1 a z_2) b[(x_1 a x_2) b (y_1 a y_2)]\} \\
 &= [z_3 a(z_1 a z_2)] b\{(x_3 b y_3) a[(x_1 a x_2) b (y_1 a y_2)]\} \\
 &= a(z_1, z_2, z_3) b[a(x_1, x_2, x_3) b a(y_1, y_2, y_3)]
 \end{aligned}$$

$$= b[a(x_1, x_2, x_3), a(y_1, y_2, y_3), a(z_1, z_2, z_3)]$$

$$\therefore a(b(x_1, y_1, z_1), b(x_2, y_2, z_2), b(x_3, y_3, z_3)) = b[a(x_1, x_2, x_3), a(y_1, y_2, y_3), a(z_1, z_2, z_3)]$$

and clearly  $1 = \Pi_1, 0 = \Pi_2$

$\therefore A$  is a Pre  $A^*$ -clone over  $S$

### Case II:

Conversely Suppose

Since  $A$  is pre  $A^*$ -clone

$\Rightarrow A$  is a pre  $A^*$ -Algebra

Define for  $a \in A, x, y, z \in S$

$$xay = a(x, y, \perp)$$

$$a(x, y, z) = za_1(xa_2y)$$

$$\begin{aligned} \text{Now } za_1(xa_2y) &= a_1(z, za_2, y, \perp) \\ &= a_1(z, a_2(x, y, \perp), \perp) \\ &= a(a_2(x, y, \perp), a_1(x, y, \perp), z) \\ &= a(a(x, y, y), a(x, y, y), z) \\ &= a(x, y, z) \\ \therefore a(x, y, z) &= za_1(xa_2y) \end{aligned}$$

$$\begin{aligned} xa_1 \perp &= a_1(x, \perp, \perp) \\ &= a_1(\perp, x, \perp) \\ &= a(x, x, \perp) \\ &= xax \end{aligned}$$

$$x1x = \Pi_1(x, x, \perp) = 2$$

$$xa \sim y = a \sim (x, x, \perp) = a(y, x, \perp) = yax$$

$$\begin{aligned} (xay)bz &= b[xay, z, \perp] \\ &= b[a(x, y, \perp), z, \perp] \\ &= b[b(x, z, \perp), b(y, z, \perp), b(\perp, z, \perp)] \\ &= a[xbz, ybz, \perp bz] \\ &= (\perp bz)a_1 [(xbz)a_2 (ybz)] \\ &= xa_2 b_2 [(\perp bz), a_1 [(ybz)]] \end{aligned}$$

$\therefore (S, \perp)$  is an "if-Then-else" structure over ' $A$ '.

**2.8 Corollary:** Let  $x, y$  be two sets and  $S = \text{Rel}(x, y)$  the set of all relations from  $X$  to  $Y$ ,  $\perp$  be the empty relation. Then  $(S, \perp)$  is an "if-Then-else" structure over Every sub  $A^*$ -algebra.  $T$  of  $T_x$ .

**Proof:** Since  $S = \text{Rel}(x, y) = \text{PF}_n(X, 2^Y)$ ,  $T_x$  can be identified with  $3^x$ ; and  $(S, \perp)$  is an "if-Then-else" structure of  $3^x$

**2.9 Theorem:** If  $S$  is an "if-then-else" structure over a Boolean algebra  $B$  and  $\perp$  is a distinguished element then  $(S, \perp)$  is an "if-then-else" structure over  $A(B)$  generated by  $B$  where  $xay = xa_1 (\perp a_2 y)$ , where  $a = (a_1, a_2)$ .

**Proof:**  $xax = xa_1 (\perp a_2 x) = \perp a_2 (xa_1 x)$   
 $= \perp a_2 x = xa_2 \perp$ .

$$xly = xl_1 (\perp l_2 y) = xl_1 y = x$$

$$\begin{aligned} xa \sim y &= xa_1 (\perp a_2 y) \\ &= x(a_1 \vee a_2) \sim (\perp a_2 y) \\ &= (\perp a_2 y)(a_2 \vee a_1)x \\ &= \perp a_2 (ya_1 x) \\ &= yax \end{aligned}$$

$$\begin{aligned} (xay)bz &= \perp b_2 [(xay)b_1 z] \\ &= [\perp a_2 (\perp a_1 \perp)] b_2 \{ [\perp a_2 (xa_1 y)] b_1 [za_2 (za_1 z)] \} \\ &= [\perp a_2 (\perp a_1 \perp)] b_2 \{ (\perp b_1 z) a_2 [(xb_1 y) b_1 (za_1 z)] \} \\ &= [\perp a_2 (\perp a_1 \perp)] b_2 \{ (\perp b_1 z) a_2 [(xb_1 z) a_1 (yb_1 z)] \} \\ &= [\perp b_2 (\perp b_2 z)] a_2 \{ (\perp a_1 \perp) b_2 [(xb_1 z) a_1 (yb_1 z)] \} \\ &= (\perp bz) a_2 \{ [\perp b_2 (xb_1 z)] a_1 [\perp b_2 (yb_1 z)] \} \\ &= (\perp bz) a_2 [(xbz)] a_1 (ybz) \\ &= (\perp bz) a_2 [(xbz)] a_1 (ybz) \\ &= xa_1 b_1 [(\perp bz) a_2 (ybz)] \\ \therefore (xay)bz &= xa_1 b_1 [(\perp bz) a_2 (ybz)] \end{aligned}$$

$$\begin{aligned} x(a * b)y &= x(a * b)_1 [\perp [x(a * b)_2 y]] \\ &= xa_1 [\perp a_1 \tilde{(b_1)} \sim y] \\ &= xa_1 \{ [\perp (\tilde{b_1}) \sim y] a_1 y \} \\ &= \{ [\perp (\tilde{b_1}) \sim y] (a_1) \sim y \} a_1 x \\ &= [\perp (\tilde{b_1}) \sim y] a_1 x \\ &= xa_1 [\perp (b_1 \vee b_2) y] \\ &= xa_1 [\perp b_2 (\perp b_1 y)] \\ &= xa_1 (\perp b_1 y) \\ \therefore x(a * b)y &= xa_1 (\perp b_1 y) \end{aligned}$$

**2.10 Lemma:** Let  $A$  is a pre  $A^*$ -algebra,  $x \in B(A)$ , then  $\text{if}_x(a, a) = a, \forall a \in A$ .

**Proof:** If  $x \in B(A)$ , then

$$\begin{aligned}\text{if}_x(a, a) &= (x \wedge a) \vee (x^{\sim} \wedge a) \\ &= (x \vee x^{\sim}) \wedge a \quad (\because x \in B(A) \Rightarrow x \vee x^{\sim} = 1) \\ &= 1 \wedge a = a.\end{aligned}$$

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